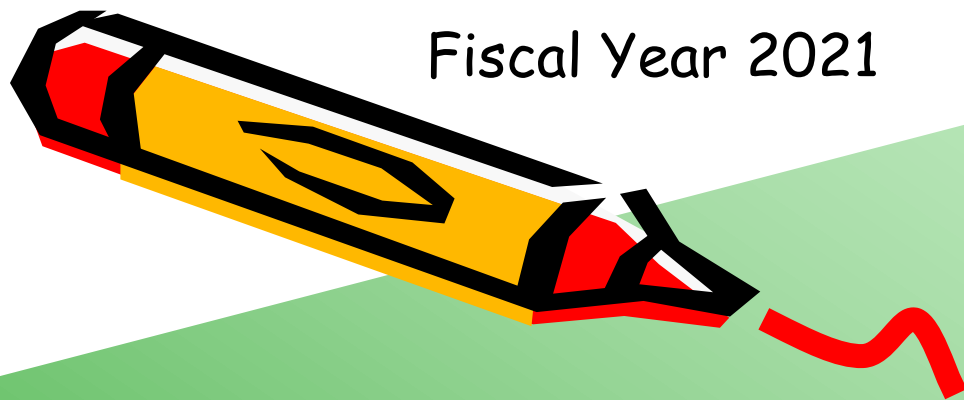


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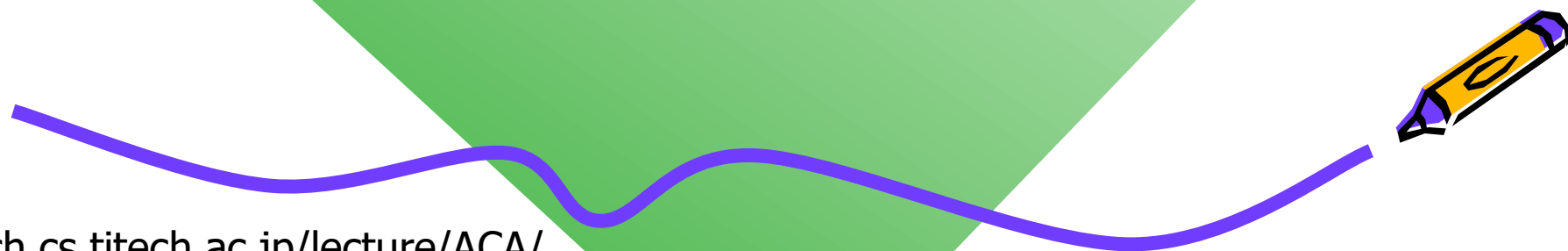
Ver. 2022-02-08a



Course number: CSC.T433  
School of Computing,  
Graduate major in Computer Science

# Advanced Computer Architecture

## 14. Final report



[www.arch.cs.titech.ac.jp/lecture/ACA/](http://www.arch.cs.titech.ac.jp/lecture/ACA/)  
Room No.W936  
Mon 14:20-16:00, Thr 14:20-16:00

Kenji Kise, Department of Computer Science  
[kise\\_at\\_c.titech.ac.jp](mailto:kise_at_c.titech.ac.jp)

# Final report



1. Submit your final report describing your answers to questions 1 - 7 in a PDF file via E-mail **by February 17, 2022**
2. **Enjoy!**



# 1. Academic paper reading

- Select an academic paper from **the list** below and
  - **In your own word**, describe the problem that the authors try to solve,
  - Describe the key idea of the proposal,
  - Describe **your opinion** why the authors could solve the problem although there may be many researchers try to solve similar problems.
- **List**
  - Prophet/critic hybrid branch prediction, ISCA'04, 2004
  - The V-Way Cache: Demand Based Associativity via Global Replacement, ISCA'05, 2005
  - Emulating Optimal Replacement with a Shepherd Cache, MICRO-40, 2008
  - A new case for the TAGE branch predictor, MICRO-44, 2011
  - Skewed Compressed Caches, MICRO-47, 2014
  - Focused Value Prediction, ISCA, 2020



## 2. MIPS assembly programming

- Write MIPS assembly code `asm1.s` for `code1.c` in C.

```
int sum = 0;
int i, j;
for (i=0; i<100; i++)
    for (j=0; j<100; j++) sum += (j+i);
```

`code1.c`

- Write MIPS assembly code `asm2.s` for `code2.c` in C.

```
int A[200];
int sum = 0;
int i;
for (i=0; i<200; i++) A[i] = i;           /* initialize the array */
for (i=1; i<200; i++) A[i] = A[i-1] + A[i]; /* compute */
for (i=0; i<200; i++) sum += A[i];        /* obtain the sum */
```

`code2.c`



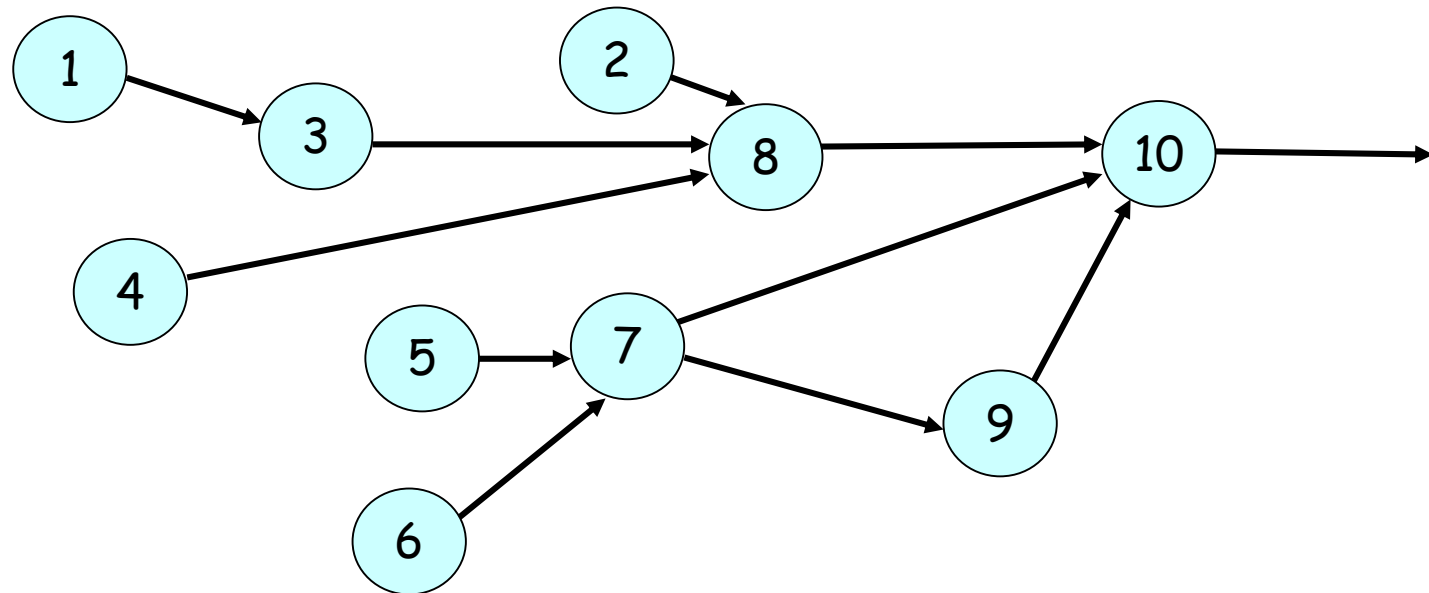
### 3. Pipelined processor

- Design a five stage pipelined scalar processor supporting MIPS `add`, `addi`, `lw`, `sw`, and `bne` instructions in Verilog HDL. Please download `proc08.v` from the support page and refer it.  
Note that you do not need to implement data forwarding.
- Verify the behavior of designed processor using `asm1.s` and `asm2.s`.  
You may insert NOP instructions if necessary.
- The report should include a block diagram, a source code in Verilog HDL, the description of the changes of the code, and obtained waveforms of your design.

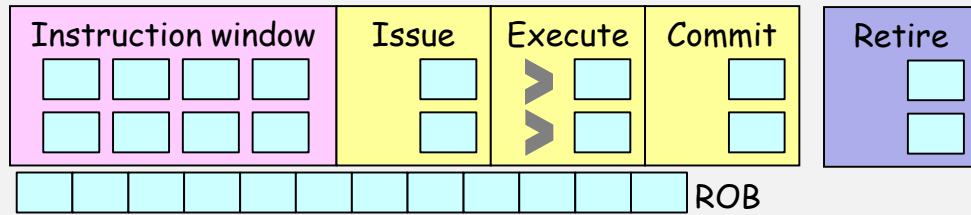


## 4. OoO execution and dynamic scheduling

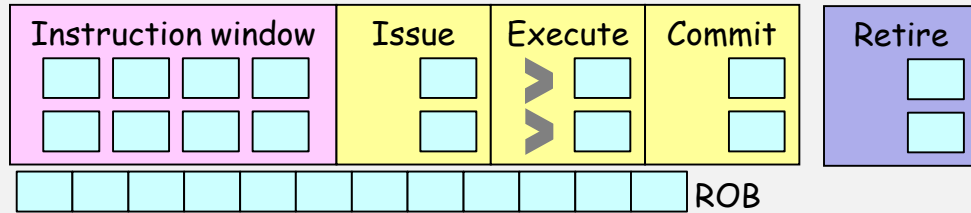
- Draw the cycle by cycle processing behavior of these 10 instructions
- Modify this dataflow graph and draw another cycle by cycle processing behavior of the graph having 10 instructions



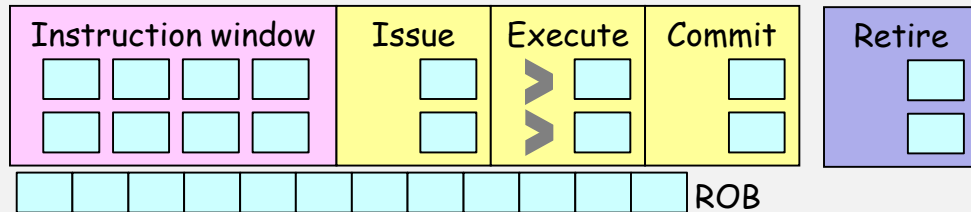
Cycle 1



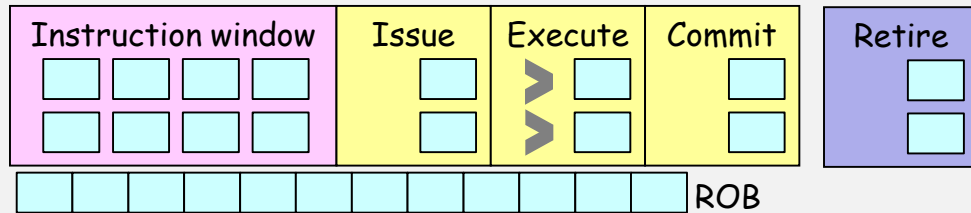
Cycle 2



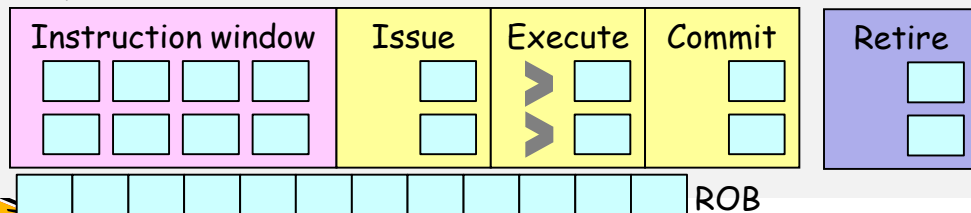
Cycle 3



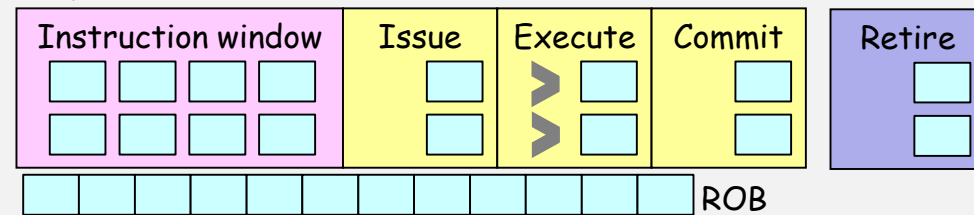
Cycle 4



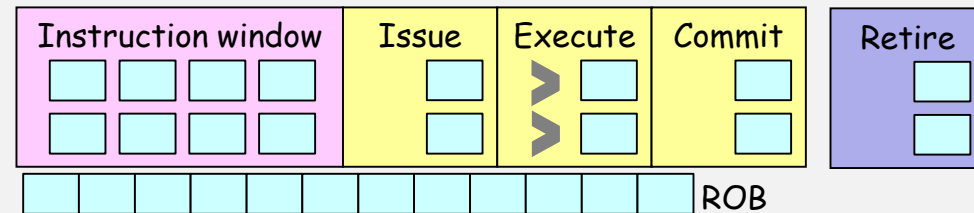
Cycle 5



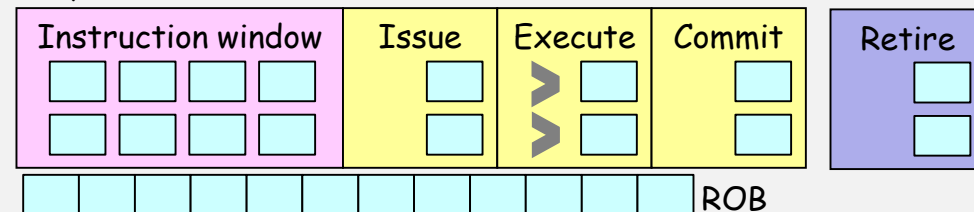
Cycle 6



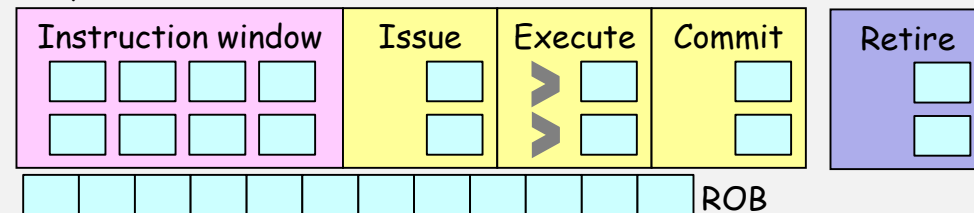
Cycle 7



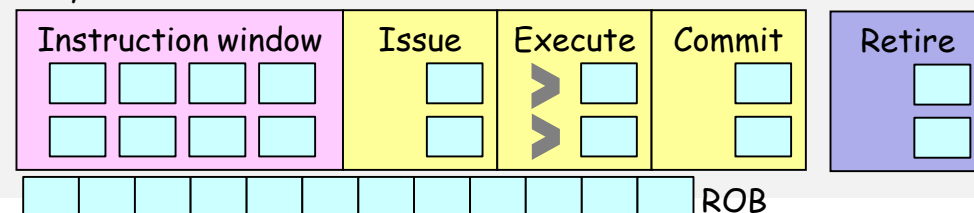
Cycle 8



Cycle 9



Cycle 10



## 5. Parallel programming (The free lunch is over)

- Describe an efficient parallel program for the following sequential program using LOCK(), UNLOCK() and BARRIER() assuming a shared memory architecture
- Explain why your code runs correctly and why your code is efficient.

```
#define N 8      /* the number of grids */
#define TOL 15.0 /* tolerance parameter */
float A[N+2], B[N+2];

void solve () {
    int i, done = 0;
    while (!done) {
        float diff = 0;
        for (i=1; i<=N; i++) { /* use A as input */
            B[i] = 0.333 * (A[i-1] + A[i] + A[i+1]);
        }
        for (i=1; i<=N; i++) { /* use B as input */
            A[i] = 0.333 * (B[i-1] + B[i] + B[i+1]);
            diff = diff + fabsf(B[i] - A[i]);
        }
        if (diff < TOL) done = 1;
        for (i=0; i<=N+1; i++) printf("%.2f ", B[i]);
        printf("| diff=%.2f\n", diff); /* for debug */
    }
}

int main() {
    int i;
    for (i=1; i<N-1; i++) A[i] = B[i] = 100+i*i;
    solve();
}
```

main02.c



## 6. Building blocks for synchronization

- **Fetch-and-increment** reads an original value from memory and increments (adds one to) it in memory atomically
- Implement fetch-and-increment (FAI) using the load-linked/store-conditional instruction pair
  - Refer the discussion of implementing an atomic exchange EXCH
- Implement BARRIER() using FAI



## 7. Cache coherence protocols



- Select your favorite commercial multi-core processor
  - Describe the memory organization including caches and main memory
    - cache line size, write policy, write allocate/no-allocate, direct-mapped/set-associative, the number of caches (L1, L2, and L3?)
  - Describe the cache coherence protocol used there

