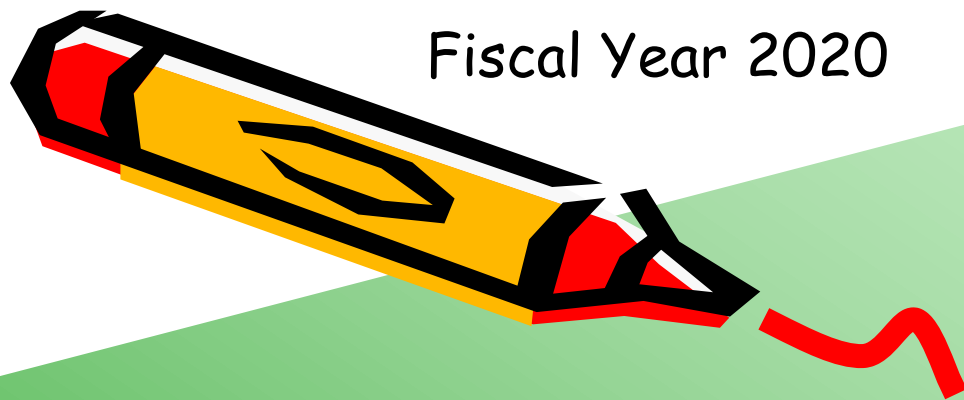


Fiscal Year 2020

Ver. 2020-12-16a



Course number: CSC.T433  
School of Computing,  
Graduate major in Computer Science

# Advanced Computer Architecture

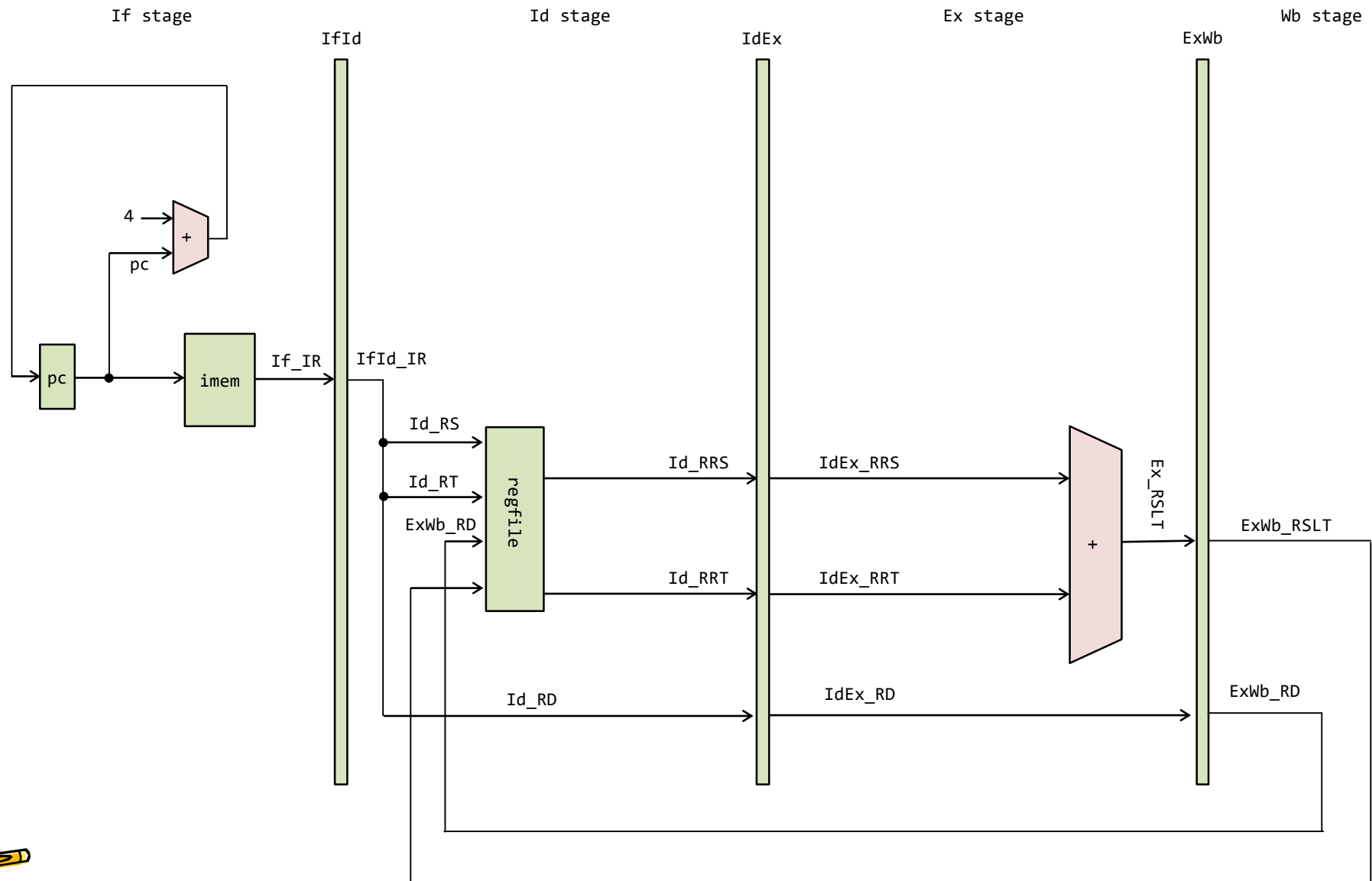
## 5. Instruction Level Parallelism: Concepts and Challenges



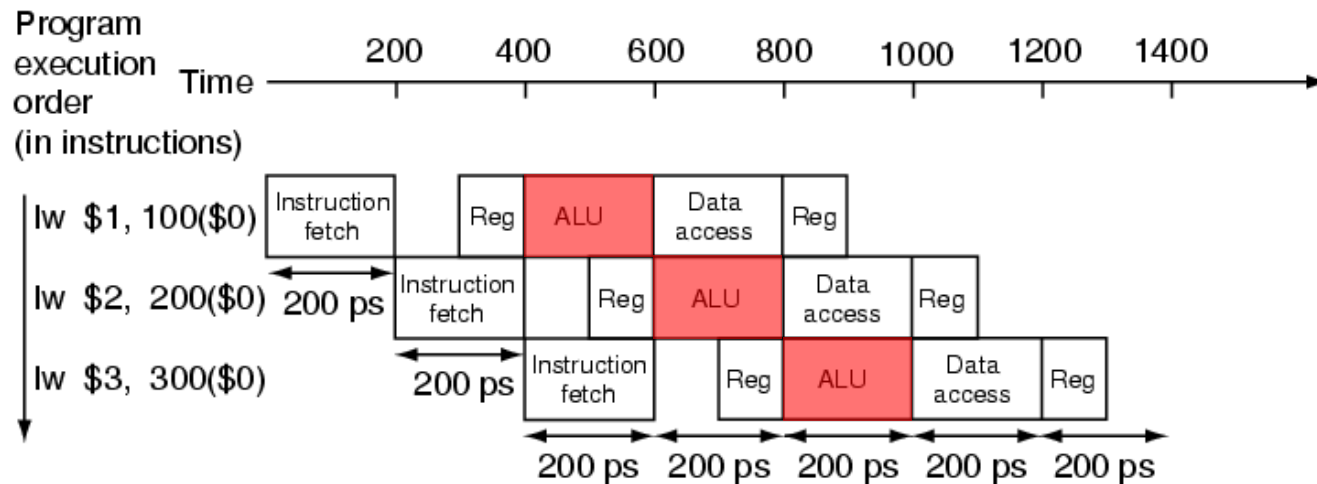
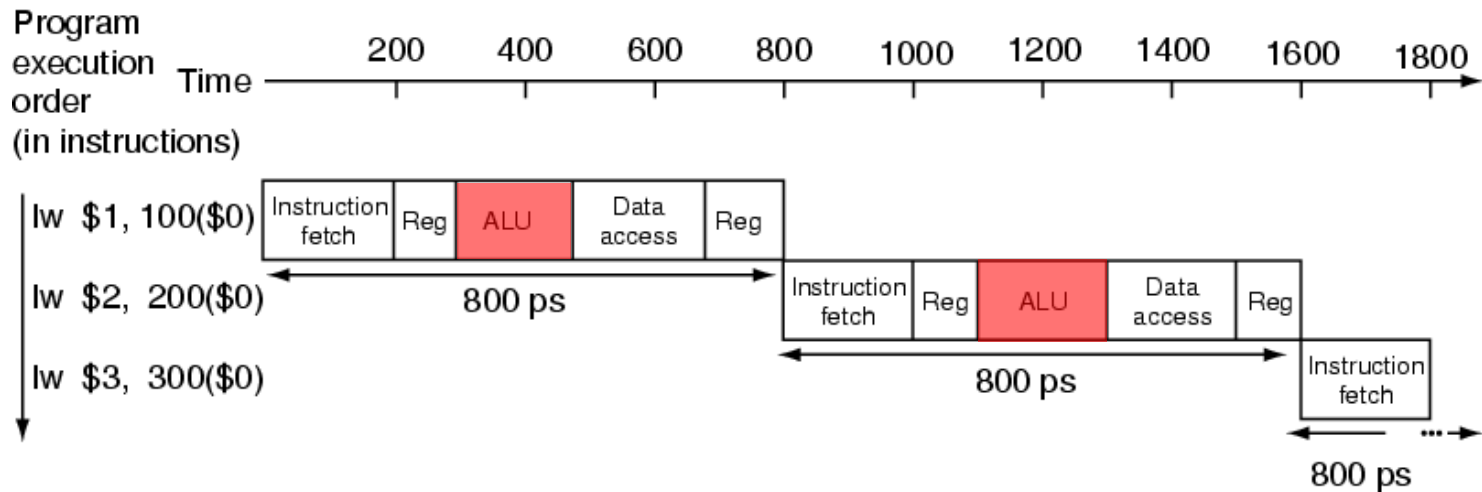
[www.arch.cs.titech.ac.jp/lecture/ACA/](http://www.arch.cs.titech.ac.jp/lecture/ACA/)  
Room No.W936  
Mon 14:20-16:00, Thr 14:20-16:00

Kenji Kise, Department of Computer Science  
[kise\\_at\\_c.titech.ac.jp](mailto:kise_at_c.titech.ac.jp)

# Four stage pipelined processor supporting ADD, which does not adopt data forwarding (proc06.v, Assignment 3)

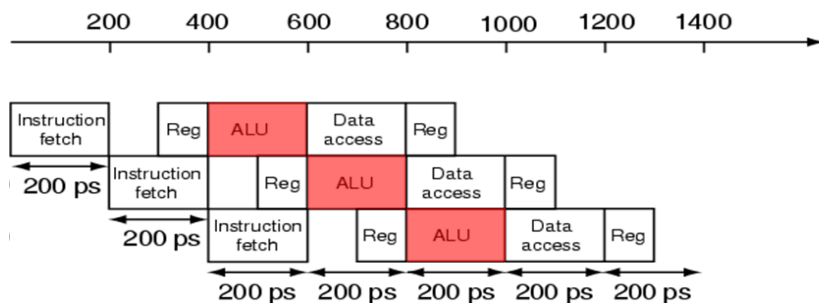


# Single-cycle and pipelined processors

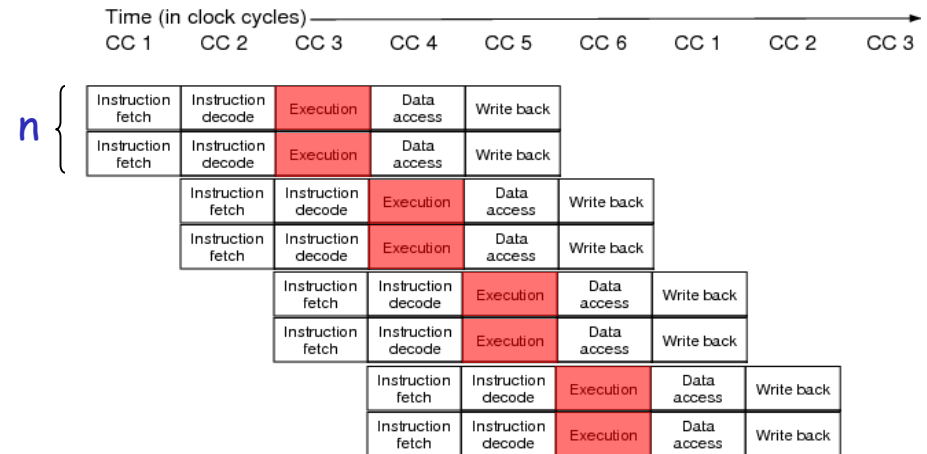


# Scalar and Superscalar processors

- **Scalar processor** can execute at most one single instruction per clock cycle using one ALU.
  - IPC (Executed Instructions Per Cycle) is less than 1.
- **Superscalar processor** can execute more than one instruction per clock cycle by executing multiple instructions using multiple pipelines.
  - IPC (Executed Instructions Per Cycle) can be more than 1.
  - using  $n$  pipelines is called  $n$ -way superscalar



(a) pipeline diagram of scalar processor



(b) pipeline diagram of 2-way superscalar processor

# Exercise: datapath of a 2-way superscalar



- Datapath of a 2-way superscalar processor supporting ADD, which does not adopt data forwarding





# Assignment 4



1. Design a four stage pipelined **2-way superscalar** processor supporting MIPS **add** instruction in Verilog HDL. Please download **proc06.v** from the support page and refer it.
2. Verify the behavior of designed processor using following assembly code  
assuming initial values of  $r[1]=22$ ,  $r[2]=33$ ,  $r[3]=44$ , and  $r[4]=55$ 
  - `add $0, $0, $0 #`
  - `add $0, $0, $0 #`
  - `add $1, $1, $1 #`
  - `add $2, $2, $2 #`
  - `add $3, $3, $3 #`
  - `add $4, $4, $4 #`
3. Submit **your report** in a PDF file via E-mail by the beginning of the next lecture.
  - The report should include a block diagram, a source code in Verilog HDL, and obtained waveforms of your design.
  - E-mail address : `report@arch.cs.titech.ac.jp`
  - E-mail title: Assignment of Advanced Computer Architecture



# Exploiting Instruction Level parallelism (ILP)

- A superscalar processor has to handle some flows efficiently to exploit ILP
  - **Control flow**
    - To execute  $n$  instructions per clock cycle, the processor has to fetch at least  $n$  instructions per cycle.
    - The main obstacles are branch instruction (BNE, BEQ)
    - Another obstacle is instruction cache
  - **Register data flow**
  - **Memory data flow**





# MIPS Control Flow Instructions

- MIPS **conditional branch** instructions:

bne \$s0, \$s1, Lbl # go to Lbl if \$s0≠\$s1

beq \$s0, \$s1, Lbl # go to Lbl if \$s0=\$s1

- Ex: **if (i==j) h = i + j;**

bne \$s0, \$s1, Lbl1

add \$s3, \$s0, \$s1

Lbl1: ...

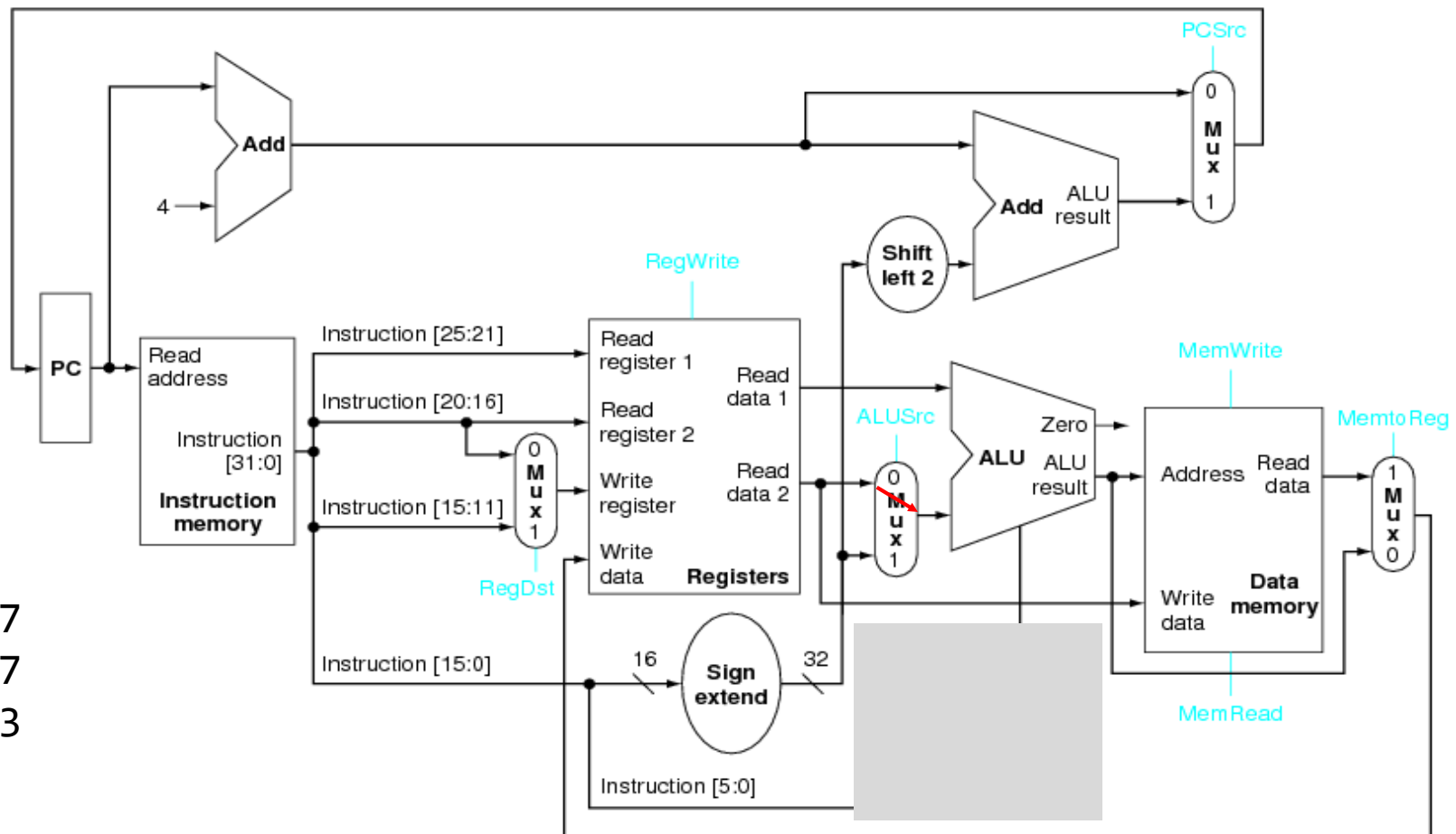
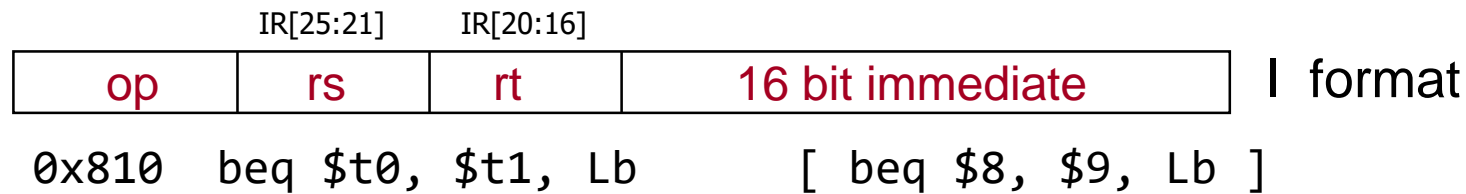
- Instruction Format (**I** format):

op	rs	rt	16 bit offset
----	----	----	---------------

- How is the branch destination address specified?



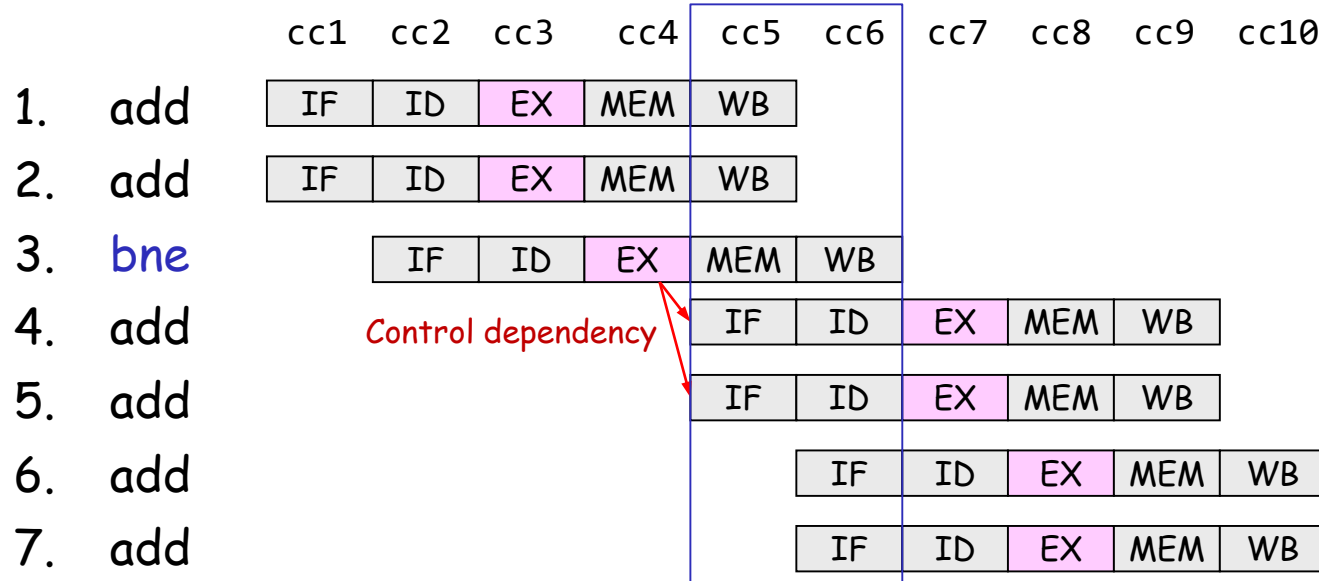
# Datapath of processor supporting **ADD, ADDI, LW, SW, BNE, BEQ**



\$8 = 7  
 \$9 = 7  
 imm = -3

# Why do branch instructions degrade IPC?

- The branch taken / untaken is determined in execution stage of the branch.
- The conservative approach of stalling instruction fetch until the branch direction is determined.

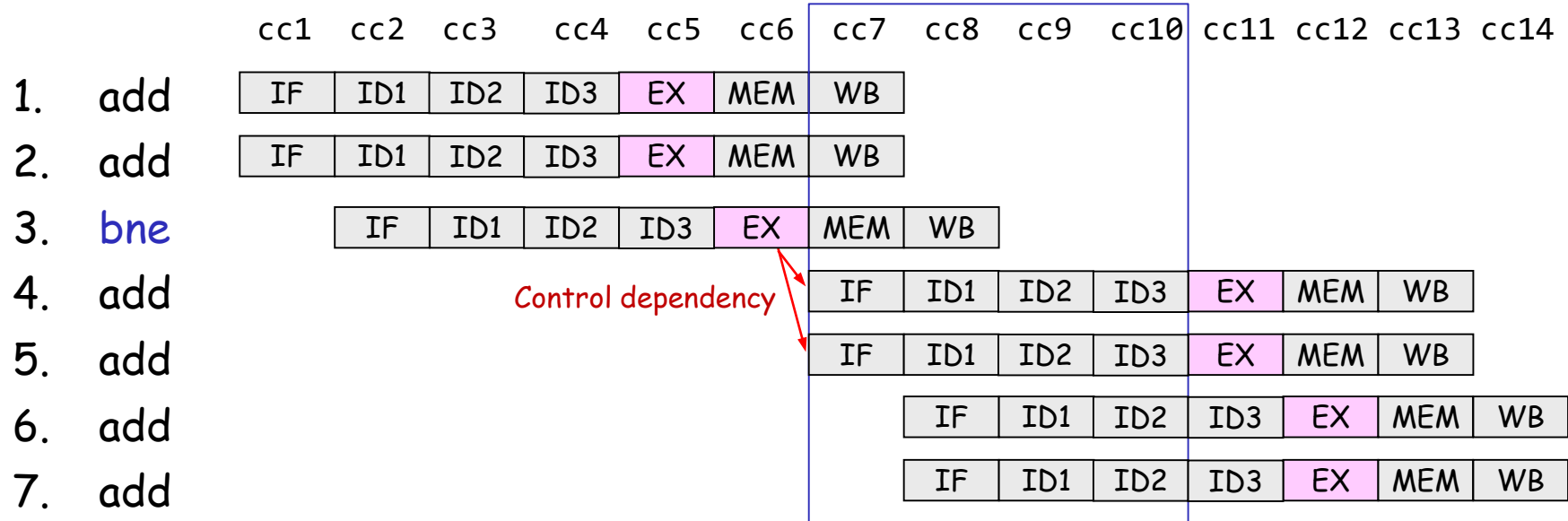


2-way superscalar processor executing instruction sequence with a branch

Note that because of a branch instruction, only one instruction is executed in cc4 and no instructions are executed in CC6 and CC7. This reduces the IPS.

# Deeper pipeline

- In conservative approach, IPC degradation will be significant by deeper pipeline



2-way superscalar adopting deeper pipeline executing instruction sequence with a branch



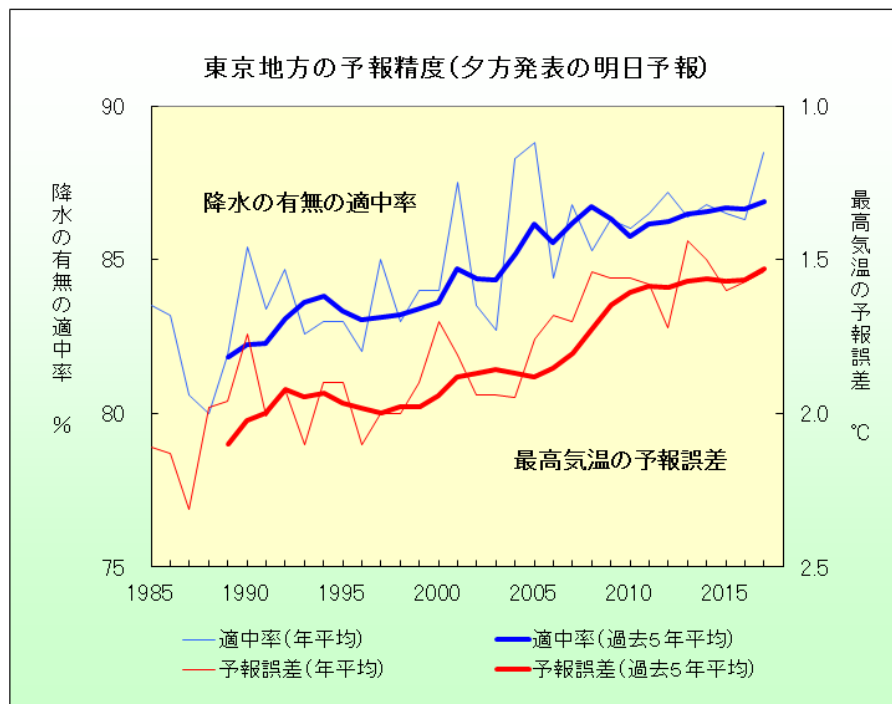
# Branch predictor



- A branch predictor is a digital circuit that tries to guess or predict which way (**taken** or **untaken**) a branch will go before this is known definitively.
  - A random predictor will achieve about a 50% hit rate because the prediction output is 1 or 0.
  - Let's guess the accuracy. What is the accuracy of typical branch predictors for high-performance commercial processors?



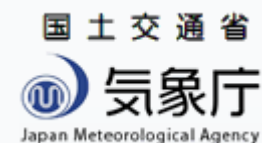
# Prediction Accuracy of weather forecasts



平成29年(2017年)までを表示しています。次の更新は平成31年(2019年)1月31日頃の予定です。

Tomorrow will be rainy?

年平均	北海道	東北	関東甲信	東海	北陸	近畿	中国	四国	九州北部	九州南部	沖縄	全国平均
明日	79	81	85	85	84	84	84	84	85	85	79	83
明後日	75	77	81	82	80	80	81	80	81	81	75	79
3日目	71	72	76	77	75	76	76	77	76	76	71	75
4日目	68	70	74	74	72	73	73	74	73	73	69	72
5日目	66	67	72	72	69	71	71	72	71	70	68	70
6日目	65	65	70	70	66	70	69	71	70	68	67	68
7日目	63	64	69	68	64	67	67	69	68	67	65	67
3~7日目平均	67	68	72	72	69	71	71	73	72	71	68	70

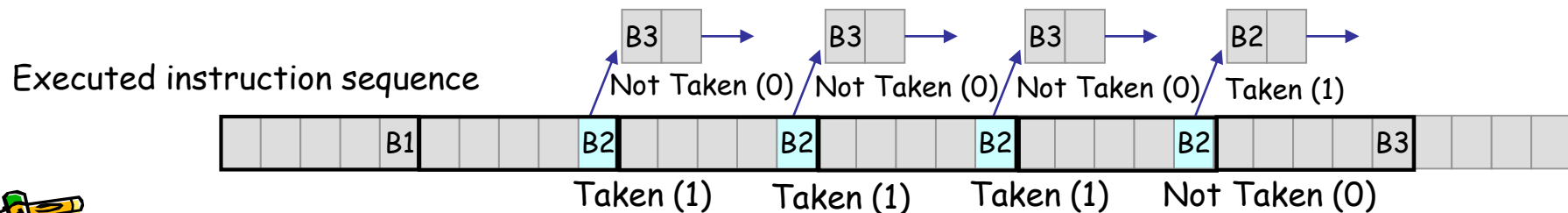
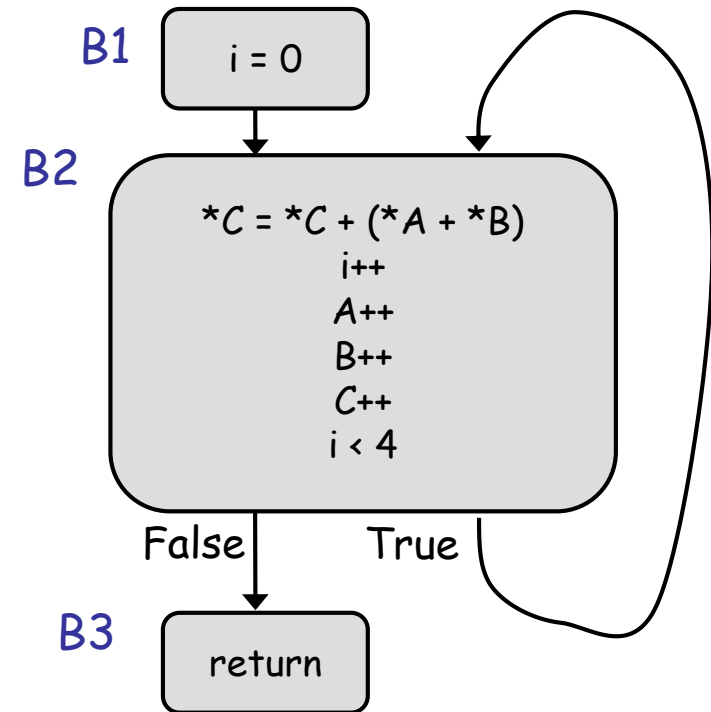


天気予報の予測精度向上に期待 - 気象庁が新スパコンを6月より稼動

マイナビニュース

# Sample program: vector add

```
#define VSIZE 4
void vadd(long *A, long *B, long *C){
    for(i=0; i<VSIZE; i++)
        C[i] += (A[i] + B[i]);
}
```





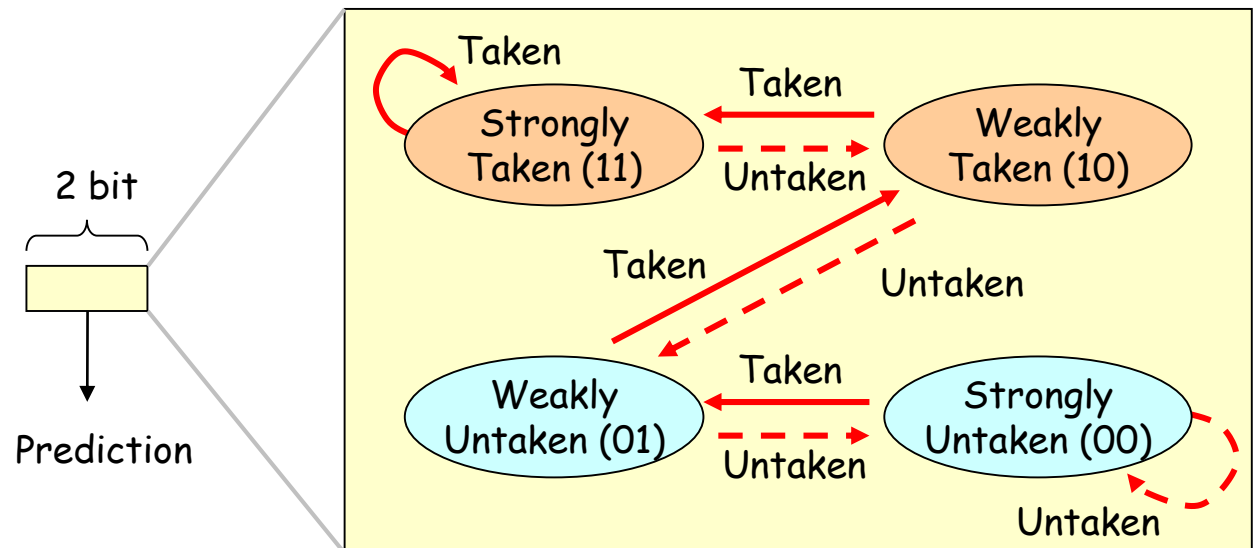
# Simple branch predictor: Branch Always

- How to predict
  - It always predicts as 1.
- How to update
  - Nothing cause it does not use any memory.



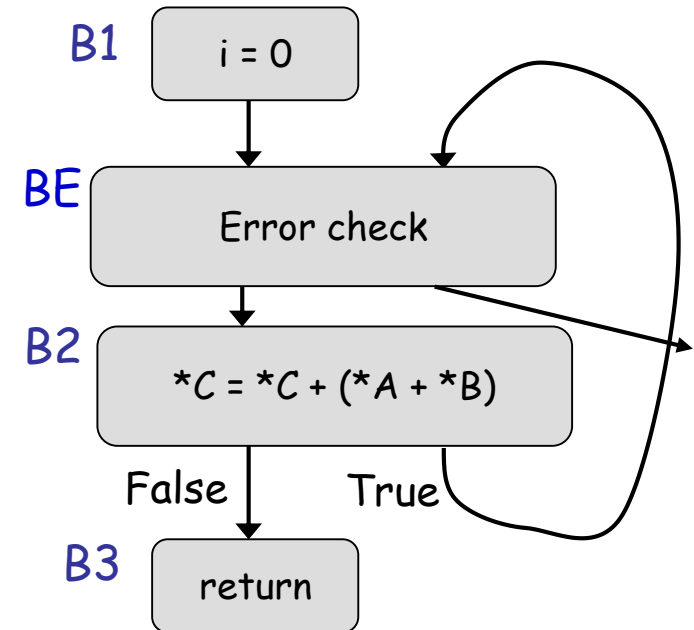
# Simple branch predictor: 2bit counter

- It uses two bit register or a counter.
- Hot to predict
  - It predicts as 1 if the MSB of the register is one, otherwise predicts as 0.
- How to update the register
  - If the branch outcome is taken and the value is not 3, then increment the register.
  - If the branch outcome is untaken and the value is not 0, then decrement the register.



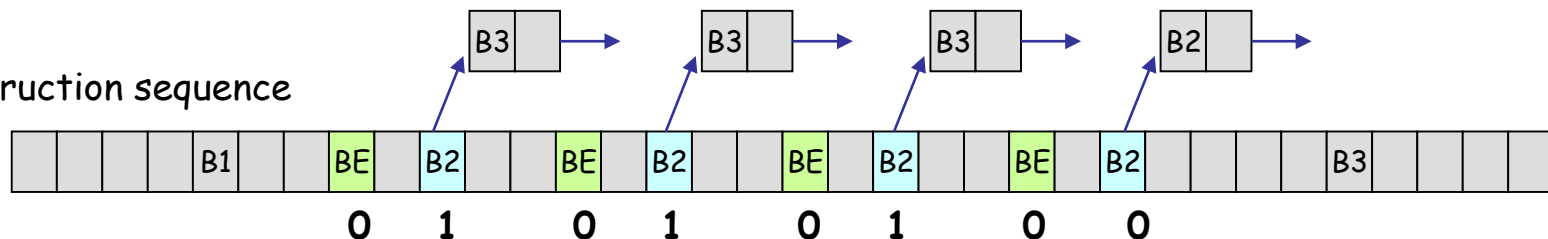
# Sample program: vector add with two branches

```
#define VSIZE 4
void vadd(long *A, long *B, long *C){
    for(i=0; i<VSIZE; i++) {
        if(A[i]<0) error_routine();
        C[i] += (A[i] + B[i]);
    }
}
```



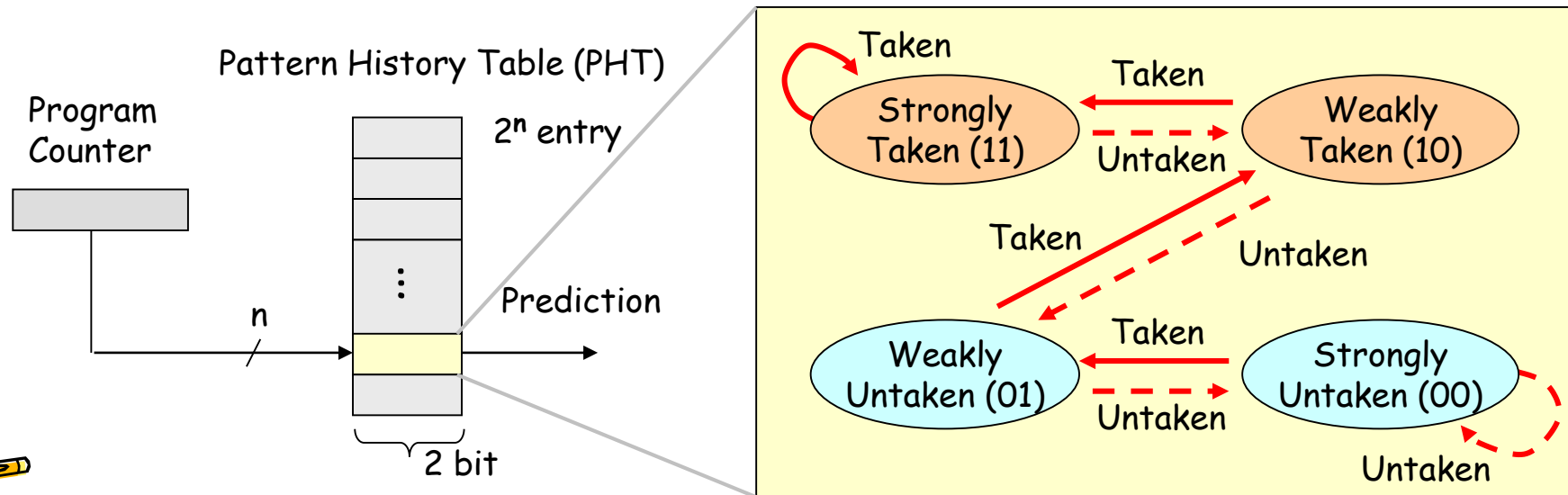
Control flow graph

Executed instruction sequence



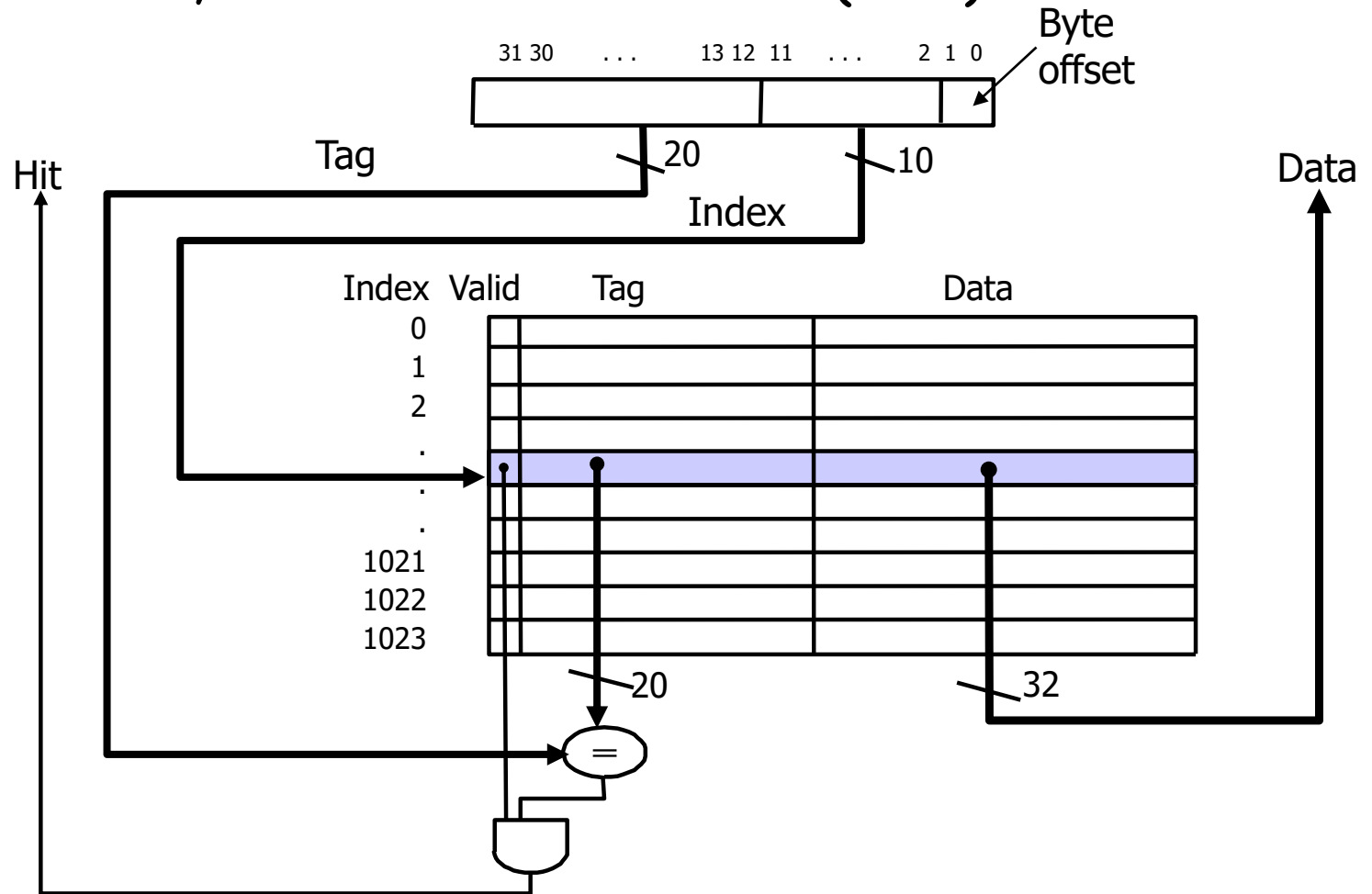
# Simple branch predictor: **bimodal**

- Program has many branch instructions. The behavior may depend on each branch. Use one counter for one branch instruction
- How to predict
  - Select one counter using PC, then it predicts 1 if the MSB of the register is one, otherwise predicts 0.
- How to update
  - Select one counter using PC, then update the counter same manner as 2bit counter.



# MIPS Direct Mapped Cache Example

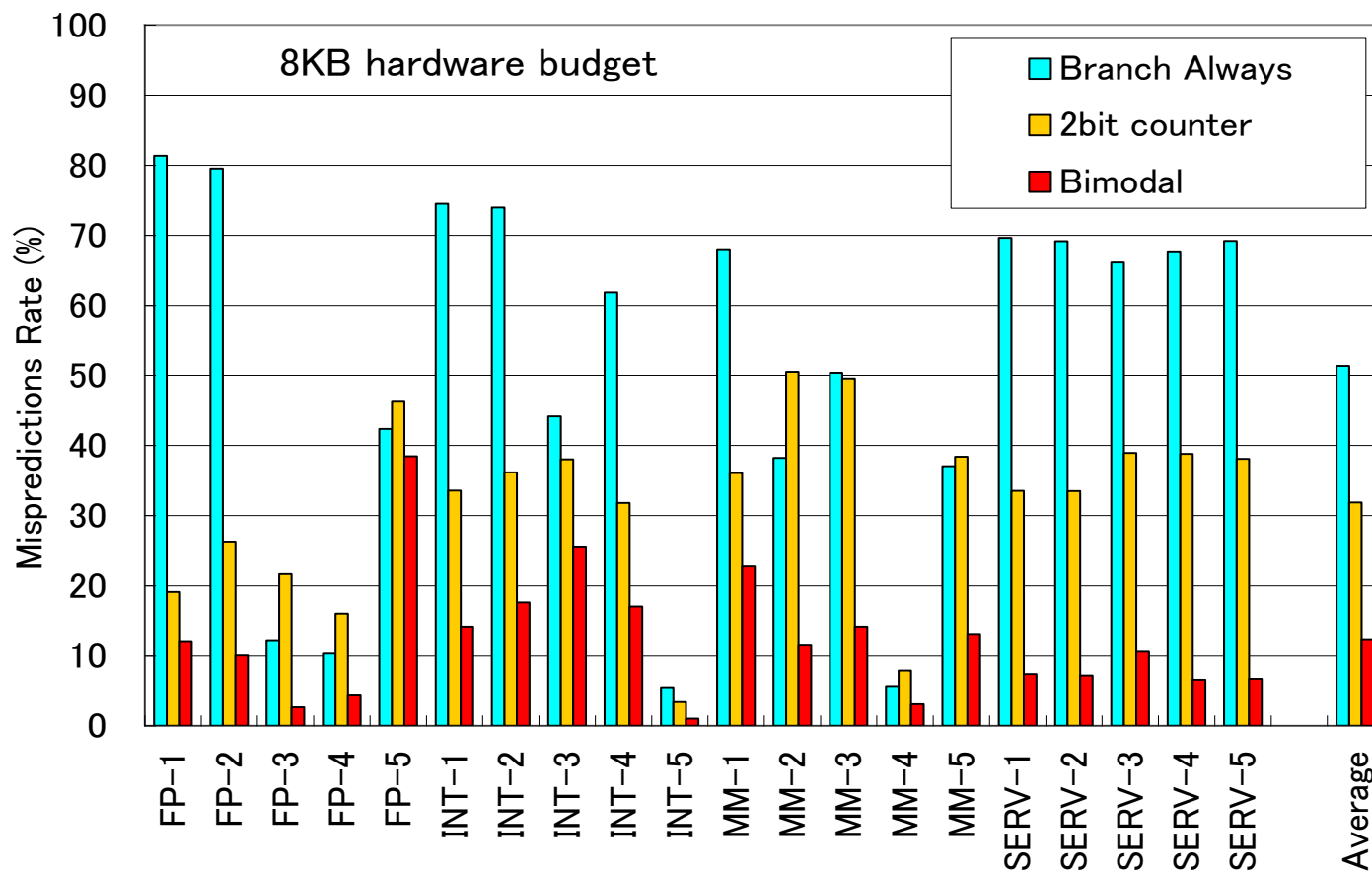
- One word/block, cache size = 1K words (4KB)



*What kind of locality are we taking advantage of?*

# Prediction accuracy of simple branch predictors

- The accuracy of branch always is about 50%.
- The accuracy of bimodal predictor of 4KB memory is about 88%.



Benchmark for CBP(2004) by Intel MRL and IEEE TC uARCH.